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Figure A1. List merge sort in memory.

```
MODULE List_merge
  SUB List_merge ( inp$( ), links( ) )
    LOCAL lower_bound, upper_bound, zero, z_times_two, i, j, k, q
    LOCAL link_01, link_02, krun, kother, inp_j$, inp_k$, vi$
    ! Check for degenerate cases of zero or one key(s)
    ! Bounds of links() used to index inp$( )
    LET lower_bound = Lbound( links )
      LET upper_bound = Ubound( links )
    IF upper_bound <= lower_bound THEN
      IF upper_bound < lower_bound THEN
        EXIT SUB
      END IF
      LET links( lower_bound ) = lower_bound
      EXIT SUB
    END IF
    ! Use "z" to represent zero, < z for end-runs, and
    ! z_times_two - p to negate a pointer p,
    ! where z = Lbound( links( ) ) - 1 and z_times_two = z*2.
    LET zero = lower_bound - 1
    LET z_times_two = 2 * zero
    ! link_01 as the second smallest; link_02 as the first smallest
    ! link_02 may be a null list
    ! link_01 is a null list if only two keys
    LET link_02 = lower_bound
    LET link_01 = link_02 + 1
    IF inp$( link_01 ) < inp$( link_02 ) THEN
      LET link_01 = link_02
      LET link_02 = link_01 + 1
    END IF
    LET krun = link_01
    LET kother = link_02
    LET inp_j$, inp_k$ = inp$( link_01 )
    FOR i = lower_bound + 2 TO upper_bound
      LET vi$ = inp$( i ) ! vi$ is current element
      IF vi$ < inp_j$ THEN ! inp_j$ is inp$(link_01)
        ! Check for new or almost minimum
        IF vi$ >= inp$( link_02 ) THEN
          LET links( i ) = link_01
          LET link_01 = i
        ELSE ! New real minimum
          LET links( i ) = links( link_02 )
          LET links( link_02 ) = link_01
          LET link_01 = link_02
          IF kother = link_02 THEN
            LET kother = i
          END IF
          LET link_02 = i
        END IF
      LET inp_j$ = inp$( link_01 )
    ELSE
      IF vi$ >= inp_k$ THEN
        ! inp_k$ is last value in run; this goes on the end
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        LET links( krun) = i
    ELSE
        ! Can not start a new run
        LET links( kother) = z_times_two - i
        LET kother = krun
    END IF
    LET krun = i
    LET inp_k$ = vi$
END IF
NEXT i
LET links( krun), links( kother) = zero
LET links( link_01) = z_times_two - links( link_01)
DO ! Make multiple passes
    LET i = links( link_01)
    LET j = links( link_02)
    IF j = zero THEN
        EXIT DO
    END IF
    LET krun = link_01
    LET kother = link_02
    DO ! Loop through runs on a single pass
        LET i = z_times_two - i
        LET j = z_times_two - j
        LET vi$ = inp$( i)
        LET inp_j$ = inp$( j)
        LET k = krun
        DO ! Merge one run
            IF vi$ <= inp_j$ THEN
                LET links( k) = i
                LET k = i
                LET i = links( i)
                IF i <= zero THEN
                    LET links( k) = j
                    DO
                        LET k = j
                        LET j = links( j)
                    LOOP WHILE j > zero
                    EXIT DO
                END IF
                LET vi$ = inp$(i)
            ELSE
                LET links( k) = j
                LET k = j
                LET j = links( j)
                IF j <= zero THEN
                    LET links( k) = i
                    DO
                        LET k = i
                        LET i = links( i)
                    LOOP WHILE i > zero
                    EXIT DO
                END IF
                LET inp_j$ = inp$(j)
            END IF
        DO
    LOOP
    ! Mark start of run

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        LET links( krun) = z_times_two - links( krun)
        LET krun = kother ! Switch output list
        LET kother = k ! Save the end of this run
    LOOP UNTIL j = zero
    LET links( krun) = i
    LET links( kother) = zero
LOOP
LET links( link_02) = link_01
LET links( link_01) = z_times_two - links( link_01)
! Complete the sorted list
LET i = link_02
FOR j = ( z_times_two - lower_bound) TO
        ( z_times_two - upper_bound) STEP - 1
    LET k = links( link_02)
    LET links( link_02) = j
    LET link_02 = k
NEXT j
FOR i = lower_bound TO upper_bound
    LET j = links( i)
    IF j < zero THEN
        LET j = z_times_two - j
        DO
            LET q = z_times_two - links( j)
            LET links( j) = k
            LET k = j
            LET j = q
        LOOP WHILE j > zero
    END IF
NEXT i
END SUB
END MODULE

```

Figure A2. Radix hash sort in memory.

```
! Number of elements to sort
LET n = 2 ^ 9
! Record length is 26 bytes
LET rec_len = 26
! Possible bit patterns in one byte = 256
LET kd = 256
! Maximum bit pattern number from zero index
LET theta = 256 - 1
!
OPTION BASE 0
DIM inp$( 0)
MAT REDIM inp$( n)
DIM inp( 0)
MAT REDIM inp( n)
MAT inp = 0
DIM hash$( 0, 0)
! Array already has element subscript 0
MAT REDIM hash$( theta, theta)
MAT hash$ = Nul$
! Initialize values in input array of first sort pass
FOR i = n TO 1 STEP - 1
    LET inp( i) = i
NEXT i
! Radix sort with hash sort for each key chunk
! Process 2-byte chunks of each record to sort
FOR h = rec_len TO 2 STEP - 2
    MAT hash$ = Nul$ ! Clear hash$ array
    ! Hash sort chunks
    FOR j = 1 TO n
        ! Must be 1..n, not n..1, to make it stable
        ! Byte order is h - 1 ... h
        ! with row as h - 1 and column as h
        LET col = ORD( inp$( inp( j)) [ h : h ] )
        LET row = ORD( inp$( inp( j)) [ h - 1 : h - 1 ] )
        LET hash$( row, col) =
            hash$( row, col) & NUM$( inp( j))
    NEXT j
    ! One greater due to subtraction operation below
    LET count = n + 1
    ! Read IEEE 8-byte record numbers into inp index
    FOR row = theta TO 0 STEP - 1
        FOR col = theta TO 0 STEP - 1
            ! Sum number of sub keys for each hash pattern
            FOR rec_num =
                LEN( hash$( row, col)) / 8 TO 1 STEP - 1
                LET count = count - 1
                LET char_num = ( rec_num - 1 ) * 8
                LET inp( count) = NUM( hash$( row, col)
                    [ char_num + 1 : char_num + 8 ] )
            NEXT rec_num
        NEXT col
    NEXT row
NEXT h
END
```

Figure A3. Performance for list merge sort and radix hash sort.

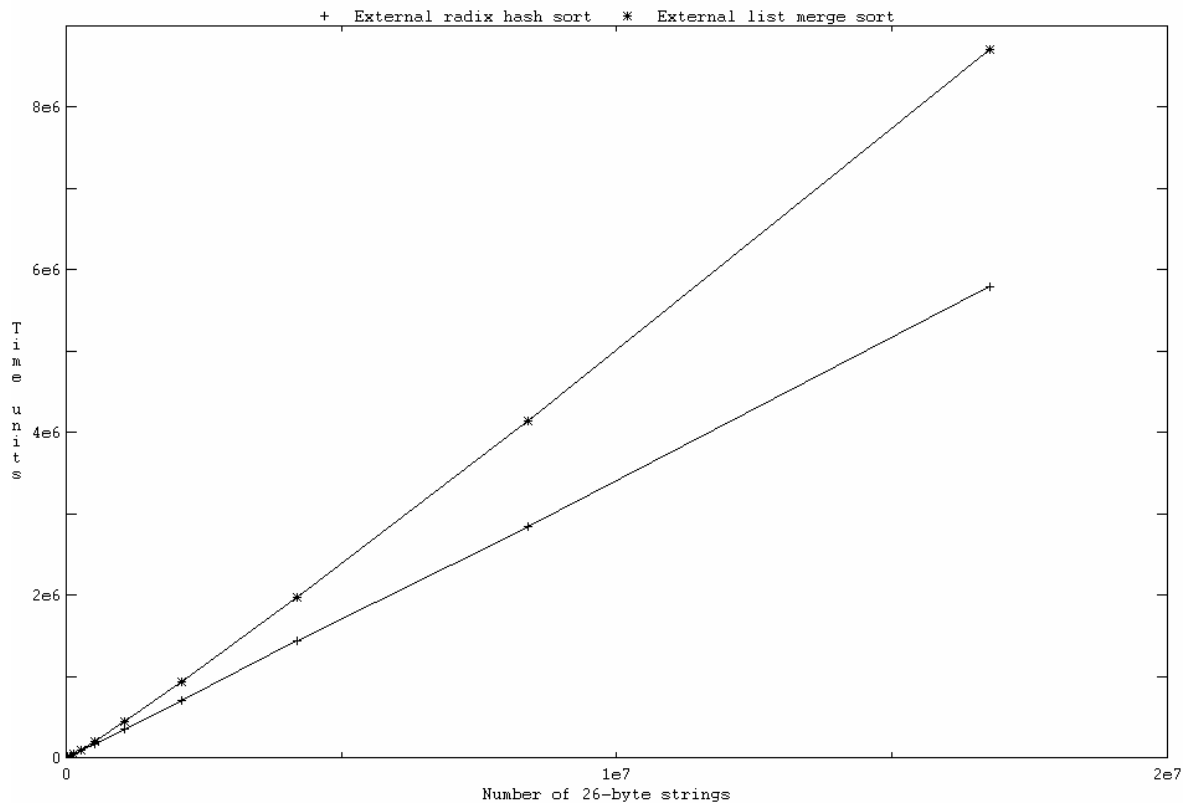


Figure A4. Logarithmic performance of list merge and radix hash sorts.

